

# Point-free Descriptive Set Theory and Algorithmic Randomness

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$\sigma$ -frames . . .

A  $\sigma$ -frame  $\mathcal{O}(X)$  is a partially-ordered set with:

- countable joins  $\bigvee$  (including the empty join  $\emptyset$ ),
- finite meets  $\wedge$  (including the empty meet  $X$ ),
- satisfying the countable distributive law:

$$U \wedge \left( \bigvee_i V_i \right) = \bigvee_i U \wedge V_i .$$

A **morphism**  $p: \mathcal{O}(Y) \rightarrow \mathcal{O}(X)$ , between  $\sigma$ -frames is a function that preserves countable joins and finite meets.

We write  $\sigma\mathbf{Frm}$  for the category of  $\sigma$ -frames.

... and  $\sigma$ -locales

A  $\sigma$ -locale  $X$  is given by a  $\sigma$ -frame  $\mathcal{O}(X)$ .

A **map**  $f: X \rightarrow Y$ , between  $\sigma$ -locales  $X, Y$ , is given by a morphism  $f^{-1}: \mathcal{O}(Y) \rightarrow \mathcal{O}(X)$  of  $\sigma$ -frames.

We write  $\sigma\mathbf{Loc}$  for the category of  $\sigma$ -locales.

(N.B.,  $\sigma\mathbf{Loc} \simeq \sigma\mathbf{Frm}^{\text{op}}$ .)

## Example $\sigma$ -frames

- $\mathcal{O}(X)$  = the lattice of open subsets of a topological space.
- $\mathcal{O}(X)$  = the lattice of Borel subsets of a topological space.
- $\mathcal{O}(X)$  = the lattice of  $\Sigma_\alpha$ -subsets of a topological space, for any ordinal  $\alpha$ .

## Full subcategories of $\sigma\mathbf{Loc}$

$\sigma\mathbf{Loc}$  is the category of maps between  $\sigma$ -locales.

- The category of continuous functions between hereditarily Lindelöf sober topological spaces.
- The category of Borel-measurable functions between standard Borel spaces.

## Complements and Boolean algebras

The **complement** (if it exists) of an element  $u$  in a distributive lattice  $P$  is the (necessarily unique) element  $\bar{u} \in P$  satisfying

$$u \wedge \bar{u} = \perp \qquad u \vee \bar{u} = \top$$

If  $p : P \rightarrow Q$  is a homomorphism of distributive lattices and  $u, \bar{u}$  are complements in  $P$  then  $p(u), p(\bar{u})$  are complements in  $Q$ .

A distributive lattice  $P$  is a Boolean algebra if and only if every element of  $P$  has a complement.

Every distributive-lattice homomorphism  $p : P \rightarrow Q$  between Boolean algebras is a Boolean-algebra homomorphism.

## $\sigma$ -Boolean algebras

For any  $\sigma$ -frame  $P$ , define a  $\sigma$ -Boolean algebra  $\mathcal{B}(P)$  and homomorphism  $i: P \longrightarrow \mathcal{B}(P)$  via the following universal property.

- for every homomorphism  $p: P \rightarrow Q$ , where  $Q$  is a  $\sigma$ -Boolean algebra, there exists a unique homomorphism  $q$  such that

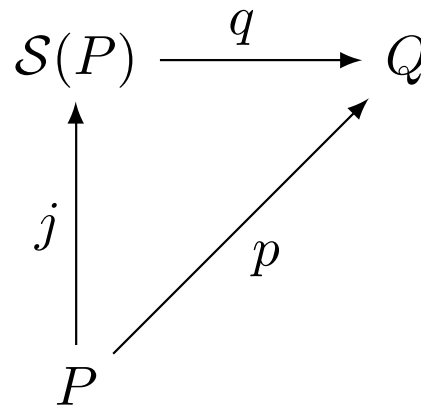
$$\begin{array}{ccc} \mathcal{B}(P) & \xrightarrow{q} & Q \\ \uparrow i & & \nearrow p \\ P & & \end{array}$$

**Theorem (CLASS).** For any quasi-Polish space  $X$ , it holds that  $\mathcal{B}(\mathcal{O}(X)) \cong \mathcal{Bor}(X)$ .

## The ‘jump’ functor

For any  $\sigma$ -frame  $P$ , define a  $\sigma$ -frame  $\mathcal{S}(P)$  and homomorphism  $j: P \longrightarrow \mathcal{S}(P)$  via the following universal property.

1. Every element in the image of  $j$  has a complement in  $\mathcal{S}(P)$ ; and
2. for every homomorphism  $p: P \rightarrow Q$ , where every element in the image of  $p$  has a complement in  $Q$ , there exists a unique homomorphism  $q$  such that





## The Borel hierarchy

**Theorem (CLASS).** For any quasi-Polish space  $X$ , it holds that  $\mathcal{S}^n(\mathcal{O}(X)) \cong \Sigma_{n+1}(X)$ .

- The classical result should generalise to ordinal-indexed iterations.
- It should hold constructively that  $\mathcal{B}$  is the free monad over the functor  $\mathcal{S}$ .
- By interpreting the definition in suitable realizability toposes, it should be possible to obtain connections with Turing degrees, the arithmetic hierarchy and the lightface hierarchy.
- ...

## Probability valuations

Write  $\overrightarrow{[0, 1]}$  for the set of ‘reals’ defined as limits of ascending sequences of rationals in  $[0, 1]$ .

A **probability** ( $\sigma$ -)valuation on a  $\sigma$ -frame  $\mathcal{O}(X)$  is a function

$$\mu : \mathcal{O}(X) \rightarrow \overrightarrow{[0, 1]}$$

satisfying

- $\mu(\emptyset) = 0$  and  $\mu(X) = 1$ .
- $\mu(u \vee v) + \mu(u \wedge v) = \mu(u) + \mu(v)$ .
- $u \leq v$  implies  $\mu(u) \leq \mu(v)$ .
- $(u_i)_i$  ascending implies  $\mu(\bigvee_i u_i) = \sup_i \mu(u_i)$ .

## The Cantor locale

Define  $\mathcal{O}(\mathbf{2}^{\mathbf{N}})$  to be the free  $\sigma$ -frame on countably many complemented generators  $(c_i)_i$ .

Intuitively, the generator  $c_i$  represents the clopen set

$$\{\alpha \in \{0, 1\}^{\omega} \mid \alpha_i = 1\}$$

**Proposition.** There is a unique probability valuation  $\lambda: \mathcal{O}(\mathbf{2}^{\mathbf{N}}) \rightarrow \overrightarrow{[0, 1]}$  such that  $\lambda(c_i) = \frac{1}{2}$  for every  $i$ .

We are endowing the Cantor ( $\sigma$ -)locale  $\mathbf{2}^{\mathbf{N}}$  with the uniform probability valuation.

## $\sigma$ -sublocales

A map  $f: X \rightarrow Y$  between  $\sigma$ -locales is said to be an **embedding** if  $f^{-1}: \mathcal{O}(Y) \rightarrow \mathcal{O}(X)$  is surjective.

The embeddings determine the notion of  **$\sigma$ -sublocale**.

The  $\sigma$ -sublocales of a  $\sigma$ -locale  $X$  are in 1–1 correspondence with congruences on  $\mathcal{O}(X)$ .

The  $\sigma$ -sublocales of  $X$  form a complete lattice  $\mathcal{S}ub(X)$  under the embedding order.

## Open $\sigma$ -sublocales

For  $v \in \mathcal{O}(X)$  define a congruence relation  $\approx_{o(v)}$  on  $\mathcal{O}(X)$  by

$$u \approx_{o(v)} u' \Leftrightarrow u \wedge v = u' \wedge v$$

It holds that

$$\mathcal{O}(X)/\approx_{o(v)} \cong \downarrow v$$

We call  $o(v)$  the **open  $\sigma$ -sublocale** determined by  $v$ .

## The $\sigma$ -locale of random sequences

For  $u, v \in \mathcal{O}(\mathbf{2}^{\mathbf{N}})$ , define:

$$u \approx v \Leftrightarrow \lambda(u) = \lambda(u \wedge v) = \lambda(v)$$

Define  $\mathcal{O}(\text{Ran}) = \mathcal{O}(\mathbf{2}^{\mathbf{N}})/\approx$ .

### Theorem

1.  $\text{Ran}$  is the intersection of all outer-measure-1  $\sigma$ -sublocales of  $\mathbf{2}^{\mathbf{N}}$ .
2.  $\text{Ran}$  is the intersection of all measure-1 open  $\sigma$ -sublocales of  $\mathbf{2}^{\mathbf{N}}$ .
3.  $\text{Ran}$  itself has outer measure 1.

(We say a  $\sigma$ -sublocale  $X \subseteq \mathbf{2}^{\mathbf{N}}$  has **outer measure** 1 if, for every open  $\sigma$ -sublocale  $X \subseteq o(u) \subseteq \mathbf{2}^{\mathbf{N}}$ , it holds that  $\lambda(u) = 1$ .)

## Points

The **terminal  $\sigma$ -locale  $\mathbf{1}$**  is given by defining  $\mathcal{O}(\mathbf{1})$  to be the free  $\sigma$ -frame on no generators.

A **point** of a  $\sigma$ -locale  $X$  is a map from the terminal  $\sigma$ -locale  $\mathbf{1}$  to  $X$ .

That is, points are given by  $\sigma$ -frame homomorphisms from  $\mathcal{O}(X)$  to  $\mathcal{O}(\mathbf{1})$ .

**Proposition**  $\mathbf{Ran}$  has no points.

## Classical points

In our intuitionistic development, there is a potentially weaker notion of **classical point** of a  $\sigma$ -locale  $X$ : a map from  $\mathbf{1}^c$  to  $X$  where

$$\mathcal{O}(\mathbf{1}^c) = \Omega_{\neg\neg} = \{p \in \Omega \mid \neg\neg p \Rightarrow p\}$$

Under classical logic,  $\mathbf{1}^c \cong \mathbf{1}$ , so classical points coincide with points.

If LEM fails, they may differ.

We can view this difference by interpreting the development in Hyland's **effective topos**  $\mathcal{E}ff$ . [Hyland 1981]



## Interpretation in $\mathcal{E}ff$

The objects in our development all produce **assemblies**.

$$|\mathcal{O}(\mathbf{2}^{\mathbf{N}})| = \{U \subseteq \{0, 1\}^{\omega} \mid U \text{ c.e. open}\}$$

$$n \mathbf{r} U \Leftrightarrow n \text{ codes a sequence } (C_i)_i \text{ of cylinders s.t. } U = \bigcup_i C_i$$

$$|\overrightarrow{[0, 1]}| = \{x \in [0, 1] \mid x \text{ left c.e.}\}$$

$$n \mathbf{r} x \Leftrightarrow n \text{ codes a sequence } (q_i)_i \text{ of rationals s.t. } x = \sup_i q_i$$

$$|\mathcal{O}(\mathbf{1})| = \{0, 1\}$$

$$n \mathbf{r} a \Leftrightarrow (a = 1 \wedge n \in K) \vee (a = 0 \wedge n \in \overline{K})$$

$$|\mathcal{O}(\mathbf{1}^c)| = \{0, 1\}$$

$$n \mathbf{r} a \Leftrightarrow \text{true}$$

## Theorem (CLASS)

1. The points of  $\mathbf{2}^{\mathbf{N}}$  in  $\mathcal{E}ff$  are in 1–1 correspondence with computable sequences  $\alpha \in \{0, 1\}^{\omega}$ .
2. The classical points of  $\mathbf{2}^{\mathbf{N}}$  in  $\mathcal{E}ff$  are in 1–1 correspondence with arbitrary sequences  $\alpha \in \{0, 1\}^{\omega}$ .
3.  $\text{Ran}$  in  $\mathcal{E}ff$  has no points.
4. The classical points of  $\text{Ran}$  in  $\mathcal{E}ff$  are in 1–1 correspondence with **Kurtz random sequences**  $\alpha \in \{0, 1\}^{\omega}$ .

A sequence  $\alpha \in \{0, 1\}^{\omega}$  is **Kurtz random** if it is contained in every measure-1 c.e. open subset  $U \subseteq \{0, 1\}^{\omega}$ . [Kurtz 1981]

## Revisiting constructive point-free descriptive set theory

There are two possible approaches to generating a  $\sigma$ -frame  $\Sigma_{\alpha+1}(X)$  from  $\Sigma_{\alpha}(X)$ .

1. Obtain  $\Sigma_{\alpha+1}(X)$  as the free  $\sigma$ -frame that adds complements to every element of  $\Sigma_{\alpha}(X)$ .

This is the ‘jump’ operation from earlier.

2. Obtain  $\Sigma_{\alpha+1}(X)$  by extending the  $\sigma$ -coframe  $\Pi_{\alpha}(X) = (\Sigma_{\alpha}(X))^{\text{op}}$  with countable joins.

Approach 1 seems the ‘correct’ approach to obtaining a rich constructive point-free descriptive set theory.

But we now follow approach 2.

## The $\sigma$ -frame $\Sigma_2(\mathbf{2}^{\mathbf{N}})$

Define  $\Sigma_2(\mathbf{2}^{\mathbf{N}})$  to be the free  $\sigma$ -frame over  $\mathcal{O}(\mathbf{2}^{\mathbf{N}})^{\text{op}}$  considered as a distributive lattice.

- There is a distributive-lattice homomorphism

$$c: \mathcal{O}(\mathbf{2}^{\mathbf{N}})^{\text{op}} \rightarrow \Sigma_2(\mathbf{2}^{\mathbf{N}})$$

- It further holds that  $c$  preserves countable meets.

We call elements of  $\Sigma_2(\mathbf{2}^{\mathbf{N}})$  in the image of  $c$  **closed**.

## The $\Sigma_2$ -reals

Define  $\overleftrightarrow{[0, 1]}$  to be the set of ‘reals’ obtained as nested sup-infs of doubly indexed sequences of rationals in  $[0, 1]$ .

**Proposition.** The probability valuation

$$\lambda: \mathcal{O}(\mathbf{2}^{\mathbf{N}})^{\text{op}} \rightarrow \overleftarrow{[0, 1]}$$

extends to a ‘probability valuation’

$$\lambda: \Sigma_2(\mathbf{2}^{\mathbf{N}}) \rightarrow \overleftrightarrow{[0, 1]}$$

Moreover,  $\lambda$  preserves countable meets of closed elements.

## Two random sub- $\Sigma_2$ -locales

Let  $\text{Ran}_1$  be the intersection of all measure-1  $\Sigma_2$  sub- $\Sigma_2$ -locales of  $\Sigma_2(\mathbf{2}^{\mathbf{N}})$ .

Let  $\text{Ran}_2$  be the intersection of all outer-measure-1 sub- $\Sigma_2$ -locales of  $\Sigma_2(\mathbf{2}^{\mathbf{N}})$ .

It is immediate that  $\text{Ran}_2 \subseteq \text{Ran}_1 \subseteq \mathbf{2}^{\mathbf{N}}$ .

**Proposition.**  $\text{Ran}_2$  (hence  $\text{Ran}_1$ ) has outer measure 1.

Theorem (CLASS).

1. The classical  $\Sigma_2$ -points of  $\text{Ran}_1$  in  $\mathcal{E}ff$  are in 1–1 correspondence with **Martin-Löf random sequences**  $\alpha \in \{0, 1\}^\omega$ .
2. The classical  $\Sigma_2$ -points of  $\text{Ran}_2$  in  $\mathcal{E}ff$  are in 1–1 correspondence with **difference random sequences**  $\alpha \in \{0, 1\}^\omega$ .

A sequence  $\alpha$  is **not ML-random** if and only if, for every confidence level  $\epsilon > 0$ , there exists (computably in  $\epsilon$ ) a c.e. open  $U_\epsilon$  with measure  $< \epsilon$  such that  $\alpha \in U_\epsilon$ . [Martin-Löf 1966]

A sequence  $\alpha$  is **difference random** if and only if it is ML-random and the halting set  $K$  is not computable relative to  $\alpha$ . [Franklin & Ng 2011]

## Some further directions

- Canonicity theorems
- Point-free and constructive measure extension theorems
- Formally develop the parallel between increasing complexity of sets and increasing complexity of real numbers